The Isar Proof Language in 2016

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Introduction

History of Isar

1999: first usable version

- primary notion of proof document (not "proof script")
- secondary notion of proof method (not "tactic")

2000–2001: various refinements

2006: minor reforms

- unfolding, obtains, literal facts: $\langle prop \rangle$
- advanced *induct* method

2016: major renovations

where "2016" means . . .

- current release: Isabelle2016 (February 2016)
- coming release: Isabelle2016-1 (November/December 2016)

Structured statements

Structured assumptions

Postfix notation for Horn-clauses:

- assume B if A_1 and A_2 for a_1 a_2
 - corresponds to assume $\bigwedge a_1 \ a_2. \ A_1 \Longrightarrow A_2 \Longrightarrow B$
 - vacuous quantifiers are omitted
- similar for **obtain**, **define**
- similar for inductive, definition, function etc.

Example: structured specifications

inductive_set star (_* [100] 100) for $R :: ('a \times 'a)$ set where

base: $(x, x) \in R \star$ for x| step: $(x, z) \in R \star$ if $(x, y) \in R$ and $(y, z) \in R \star$ for x y z

function $gcd :: nat \Rightarrow nat \Rightarrow nat$ where

$$\begin{array}{l} gcd \ x \ 0 = x \\ | \ gcd \ 0 \ y = y \\ | \ gcd \ (Suc \ x) \ (Suc \ y) = gcd \ (Suc \ x) \ (y - x) \ \text{if} \ x < y \\ | \ gcd \ (Suc \ x) \ (Suc \ y) = gcd \ (x - y) \ (Suc \ y) \ \text{if} \ \neg \ x < y \end{array}$$

Structured conclusions (goals)

Notation for Isar "eigen-context":

- premises: have B if A_1 A_2
- parameters: have B for $a_1 a_2$
- corresponds to { fix $a_1 a_2$ assume that: $A_1 A_2$ have B }
- analogous to lemma fixes $a_1 a_2$ assumes that: $A_1 A_2$ shows B

Example: Natural Deduction with structured conclusions

- conjunction introduction: have $A \land B$ if A and B
- existential introduction:
 have ∃ x. B x if B a for a
- disjunction elimination: from $\langle A \lor B \rangle$ have C if $A \implies C$ and $B \implies C$ for C
- existential elimination:

from $(\exists x. B x)$ have C if $\bigwedge x. B x \implies C$ for C

Elimination statements

```
consider \overline{x} where \overline{A} \ \overline{x} | \overline{y} where \overline{B} \ \overline{y} | \ldots \equiv
have thesis
if \bigwedge \overline{x} . \ \overline{A} \ \overline{x} \Longrightarrow thesis
and \bigwedge \overline{y} . \ \overline{B} \ \overline{y} \Longrightarrow thesis
for thesis
```

Examples:

- existential elimination:
 from ⟨∃ x. B x⟩ consider x where B x
- conjunction elimination: from $\langle A \land B \rangle$ consider A and B
- disjunction elimination:
 from ⟨A ∨ B⟩ consider A | B

Elimination and cases

- method "cases" detects its rule from chained facts
- command "case" allows name and attribute specification

Example:

```
consider x where A x | y where B y \langle proof \rangle
then have something
proof cases
  case prems: 1
  show ?thesis using prems \langle proof \rangle
  next
  case prems: 2
  show ?thesis using prems \langle proof \rangle
  qed
```

Obtain

obtain \overline{x} where $\overline{A} \ \overline{x} \langle proof \rangle \equiv$ consider \overline{x} where $\overline{A} \ \overline{x} \langle proof \rangle$ fix \overline{x} assume^{*} $\overline{A} \ \overline{x}$

- old meaning is unchanged, but foundation simplified
- is patterns now supported (with λ -lifting over the parameters)
- if / for notation available as well

Define

```
define c where c \ \overline{x} = t for \overline{x} \equiv def \ c \equiv \lambda \overline{x}. \ t
```

- syntax like **obtain**
- analogous to **definition** (e.g. object-logic equalities)
- old **def** is declared legacy

Strong vs. weak premises

- strong premises (cf. assume): show B if A_1 and A_2
- weak premises (cf. presume): show B when A_1 and A_2
- show $A_1 \Longrightarrow A_2 \Longrightarrow B$ becomes free for re-interpretation:

```
have \langle A \longrightarrow B \rangle

proof

show \langle B \rangle if \langle A \rangle \langle proof \rangle — strong premise (new in 2016)

qed

have \langle A \longrightarrow B \rangle

proof

show \langle A \Longrightarrow B \rangle \langle proof \rangle — strong premise (changed in 2016)

qed
```

Proof structure

Simplified block structure

Nesting levels:

- + goal statement (have, show etc.)
- = backwards refinement (**using**, **apply**, **supply** etc.)
- + proof
- + {
- }
- -- qed

Some consequences:

- cases in proof methods no longer special (regular context update)
- Eisbach: *match* method can use generic context for bookkeeping
- Isabelle/jEdit: clarified text folding and indentation

Structured backwards refinement

 $\langle goal \rangle$ subgoal premises *prems* for $x_1 x_2 \dots$ $\langle proof \rangle$

Example: structured apply "scripts"

```
\langle goal \rangle
subgoal by method_1
subgoal by method_2
done
```

$\langle goal \rangle$

subgoal premises *prems* for $x \ y$ using *prems* $\langle proof_1 \rangle$ subgoal premises *prems* for $u \ v \ w$ using *prems* $\langle proof_2 \rangle$ done

Proof method facts

Used facts of method expression:

- get via dynamic fact *method_facts* (useful for Eisbach method definitions)
- set via method *use*, e.g.
 - $(use \ldots in simp)$
 - $(use \ldots in \langle simp \ add : \ldots \rangle)$
- special fact *nothing* may help in odd situations

Example:

have $a: A \langle proof \rangle$ have B by (rule r) (use a in auto)

Isar document language

Document structure

Markup

- section headings (6 levels like in HTML): chapter, section, subsection, . . . , subparagraph
- text blocks: text, txt, text_raw
- uncontrolled <a>EX macros (rare)

Markdown

• implicit paragraphs and lists: itemize, enumerate, description

Document antiquotations

full form: @{name [options] arguments ...} short form:

- 1. cartouche argument: **\<^***name>*(*argument*)
- 2. no argument: \<^*name*>
- 3. standard name: (*argument*)

Notable examples:

- *cartouche*, *theory_text*: self-presentation of lsar
- *bold*, *emph*, *verbatim*, *footnote*: text styles (with proper nesting)
- noindent, smallskip, medskip, bigskip: spacing
- *cite*: formal BibT_EX items
- *path, file, dir, url, doc*: system resources

Conclusions

Conclusions

- raw proof blocks { ... } are rarely required: superseded by structured conclusions with if / for eigen-context
- big-bang integration of several cases "by *blast*" is obsolete: superseded by **consider** and enhanced method *cases*
- integration of unstructured **apply**-scripts into structured proofs works better (notably via **subgoal**)
- auxiliary method *insert* is mostly obsolete: superseded by (*use* ... **in** *method*)
- \bullet Isar document reforms make it very easy to make presentations on the spot, without $\ensuremath{\text{ETEX}}$ tinkering

TODO

- error-recovery according to block structure of proof document
- re-unify *where* / *of* in Eisbach vs. Pure
- re-unify *atomize* and *atomize_elim* as *compact*
- $\bullet \mbox{ make } compact \mbox{ the default of automated methods}$
- re-unify of *induct / induction* and *coinduct / coinduction*
- eliminate *induct_tac*, *case_tac*, *rule_tac* eventually
- proper HTML document output (e.g. presentations with **reveal.js**)
- interactive document preparation (with PIDE)